

ABFT in Dense Linear Algebra

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Motivation

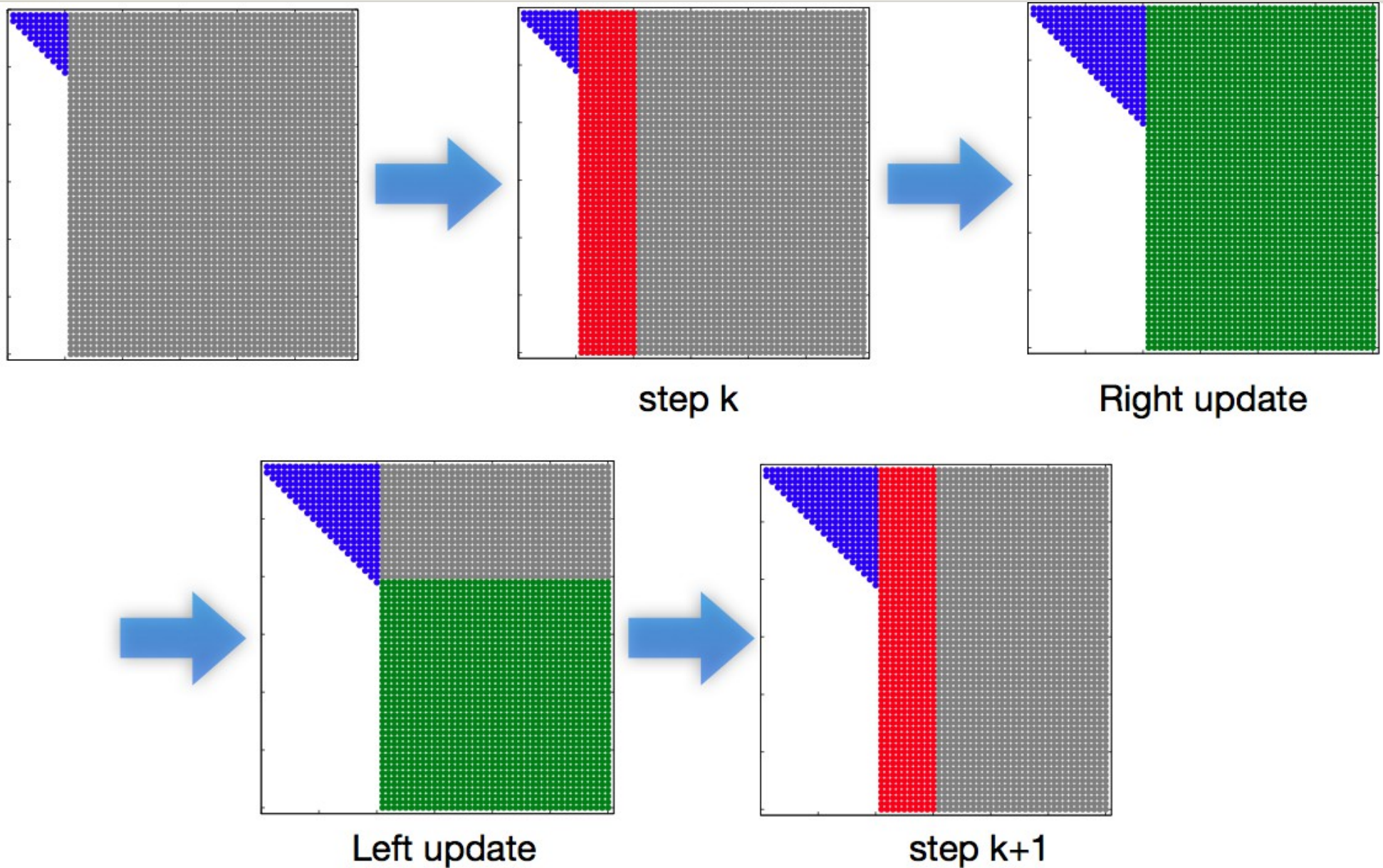
- Soft errors...
 - caused by: cosmic rays (alpha particle, high energy and/or thermal neutrons)
 - occur in practice
 - Commercial study in 2000 by Sun Microsystems
 - ASC Q supercomputer at Los Alamos in 2003
 - Jaguar (Cray XT5) at ORNL
 - Nearly 225k cores
 - 1253 separate node crashes during 537 days (Aug 2008-Feb 2010)
 - Or 2.33 failures per day Or less 10 hours of failure-free operation
 - ... and any non-ECC machine
- Accelerators are common
 - In many shared-memory systems
 - Supercomputers
 - Tianhe-1A, Titan (Cray XK7, 560k cores), Tianhe-2 (3M+ cores)
- And at Exascale ~ 1 billion threads and $MTTF < 1$ day!

Importance of a Single Node

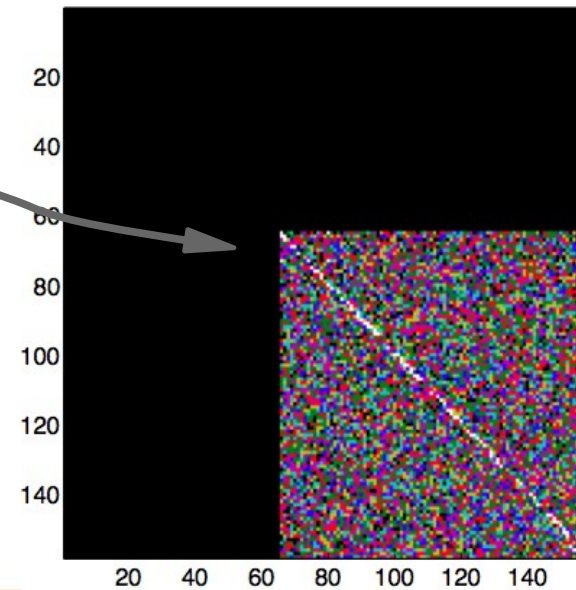
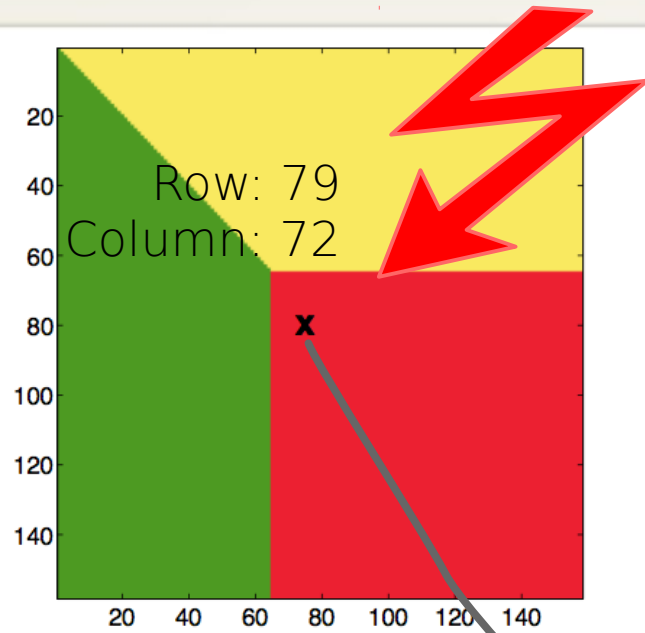
- MTTF of the entire machine depends on reliability of each node
- The MTTF of the entire machine can be statistically computed based on single-node reliability for a number of distributions
 - $\text{Exp}(1/100)$
 - $\text{Weibull}(0.7, 1/100)$
 - $\text{Weibull}(0.5, 1/100)$
- See Yves Robert's work for detailed analysis

One node $\sim 10^3$ cores	MTTF = 1 year	MTTF = 10 years	MTTF = 120 years
↓	↓	↓	↓
Exascale machine $\sim 10^6$ nodes	30 seconds	5 minutes	1 hour

Two-Sided Matrix Factorization



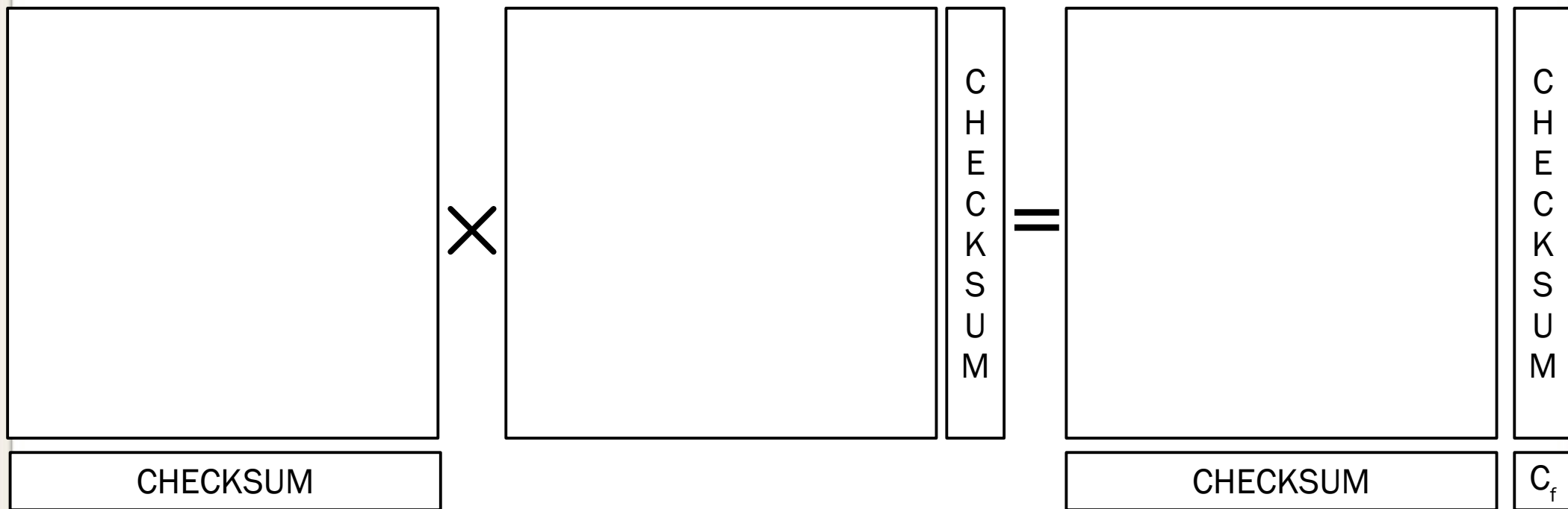
Error Propagation



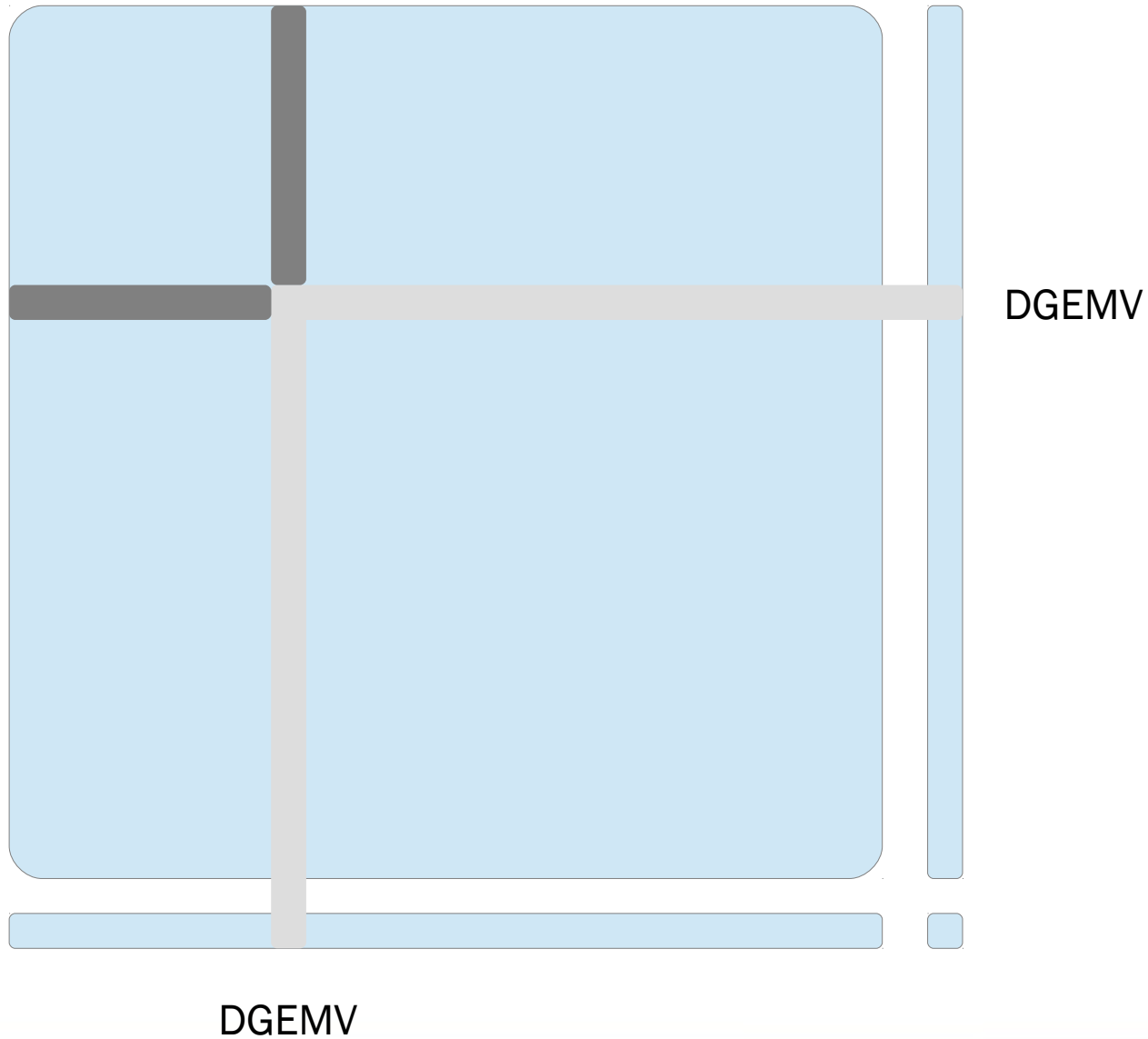
Techniques for Error Protection and Failure Recovery

- Algorithm-Based Fault Tolerance
 - Kuang-Hua Hua, Jacob Abraham, ABFT for Matrix Operations
 - Implementation on systolic arrays
 - Takes advantage of additional mathematical relationship(s)
 - Already present in algorithm
 - Introduced (cheaply, if possible) by ABFT – usually weighted sums
- Diskless checkpointing
 - Additional (small) data is kept in live processes
 - No need for full I/O checkpointing

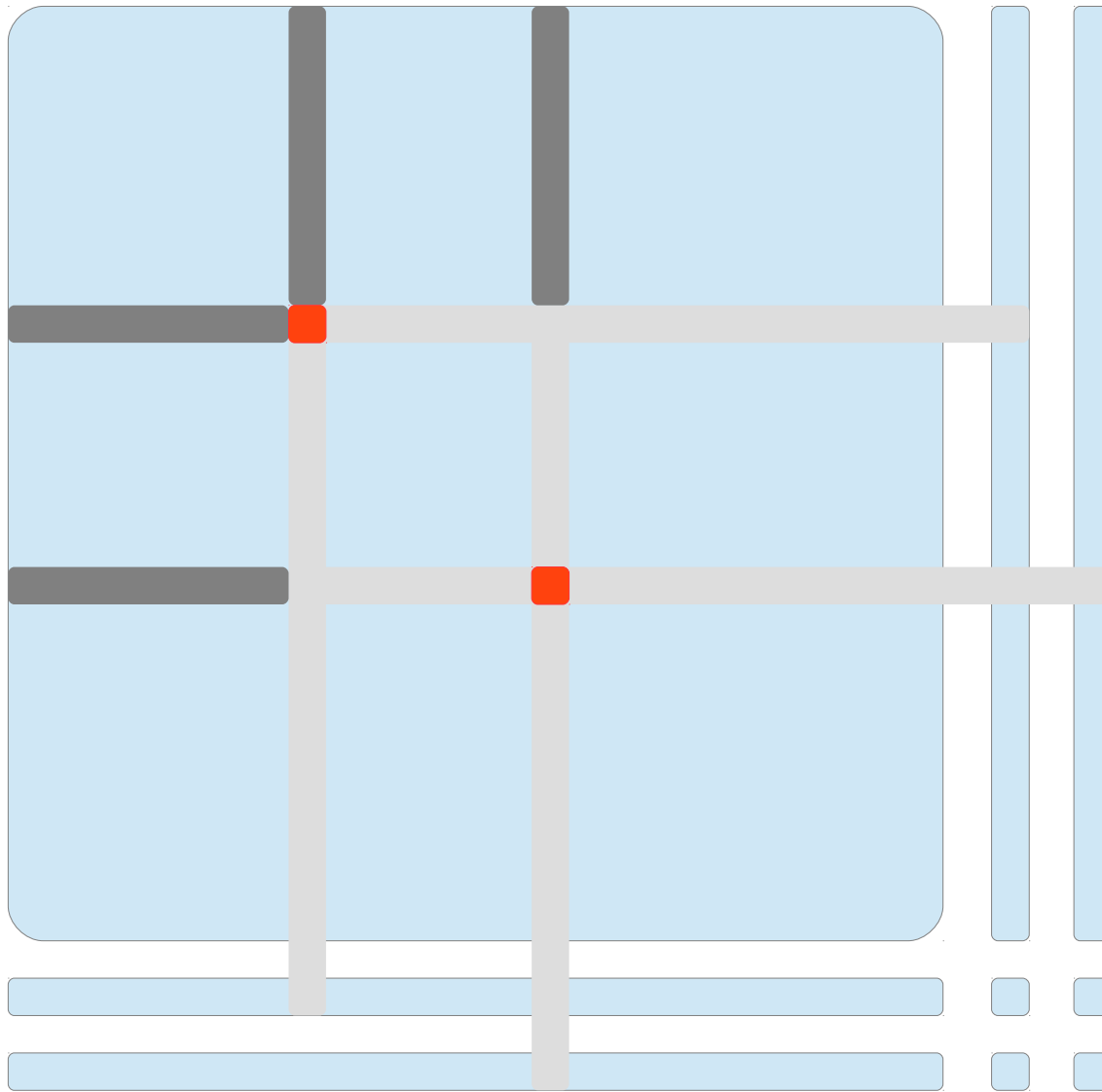
Huang&Abraham: Checksum M-M-mul



Extra Computation for Error Protection



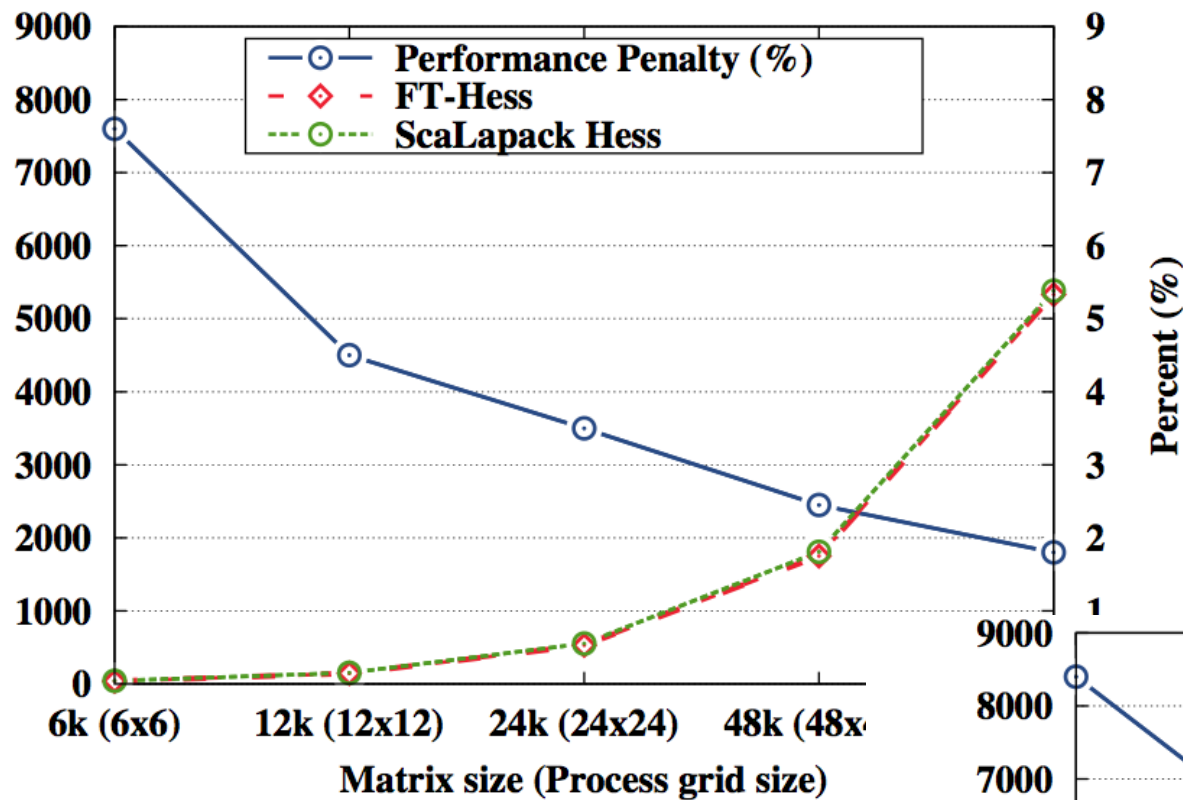
Computation for Two-Error Protection



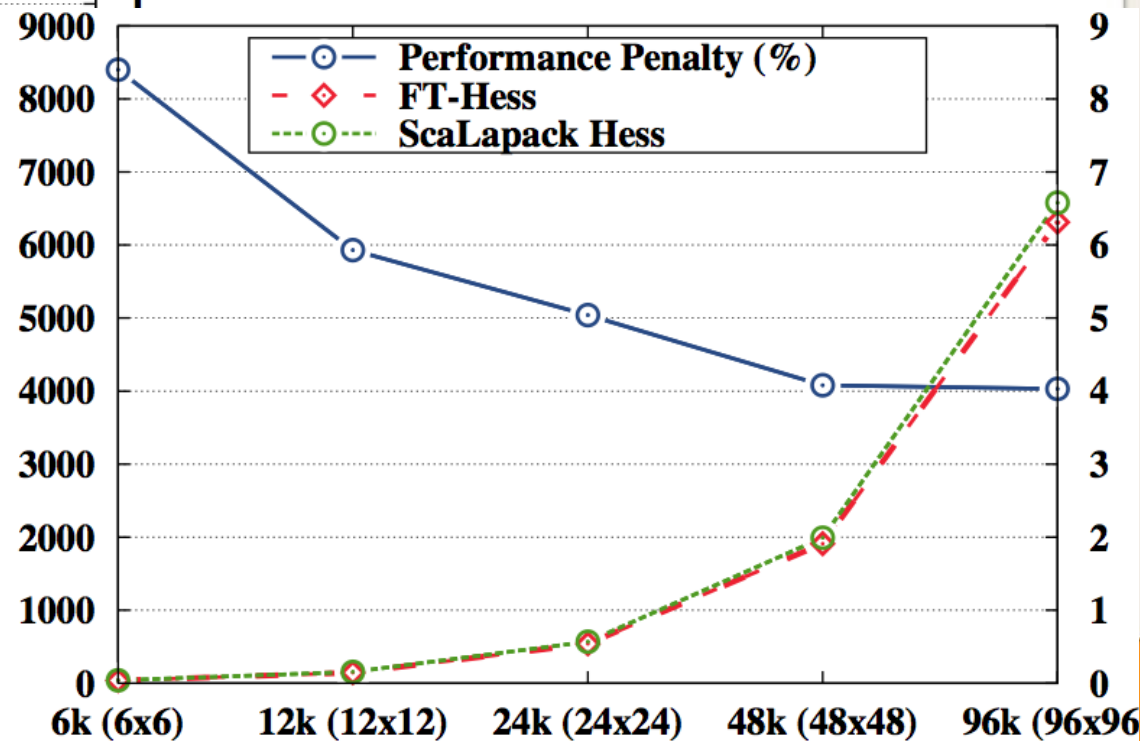
DGETRF (solve)

Experiment: Hessenberg Reduction

- ORNL Titan: Cray XK7
- AMD Opteron 6274
- 299,008 core
- Cray Gemini 3D Torus



1 error →



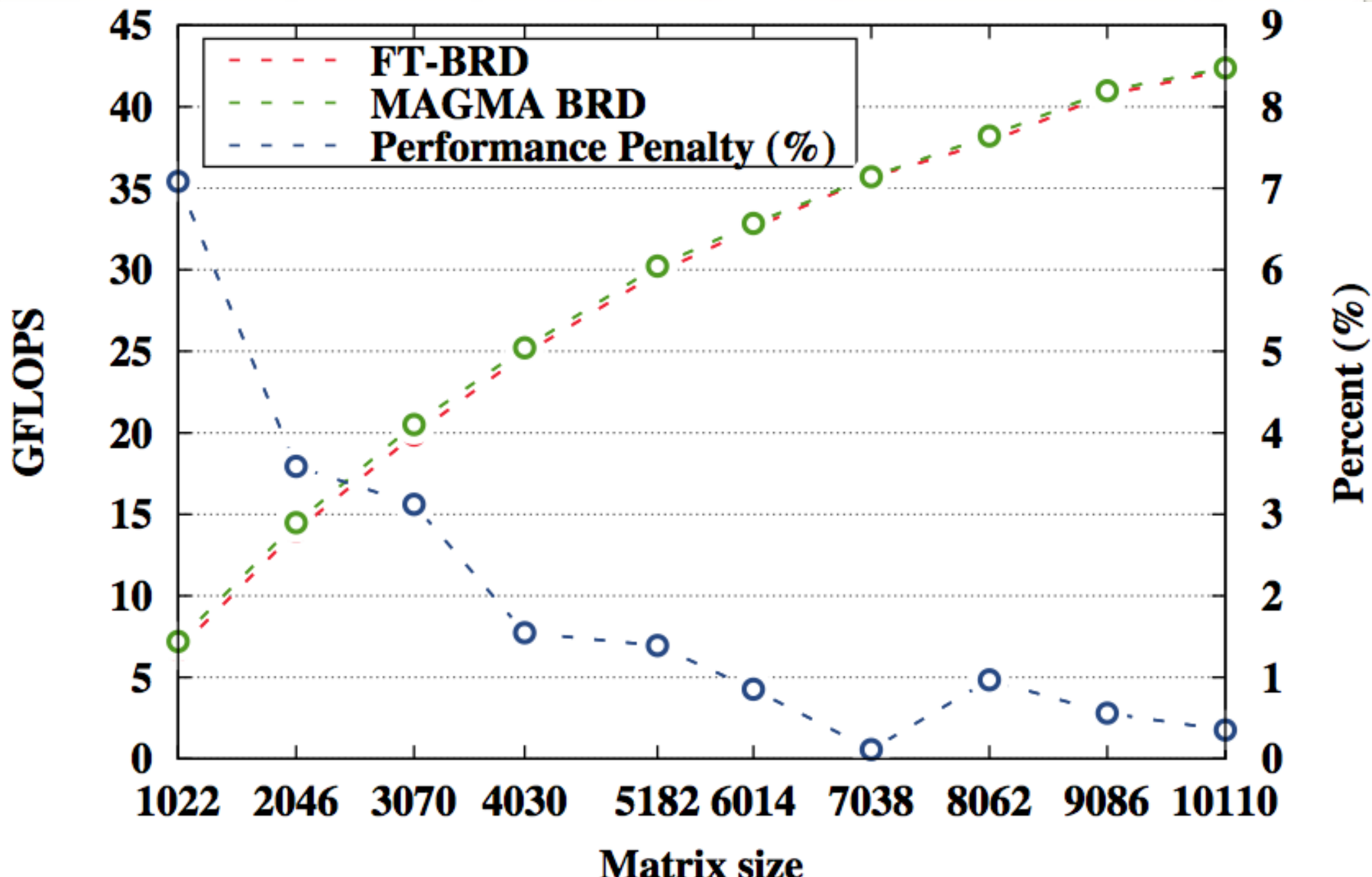
← No errors

Hessenberg Reduction: Recovery Error

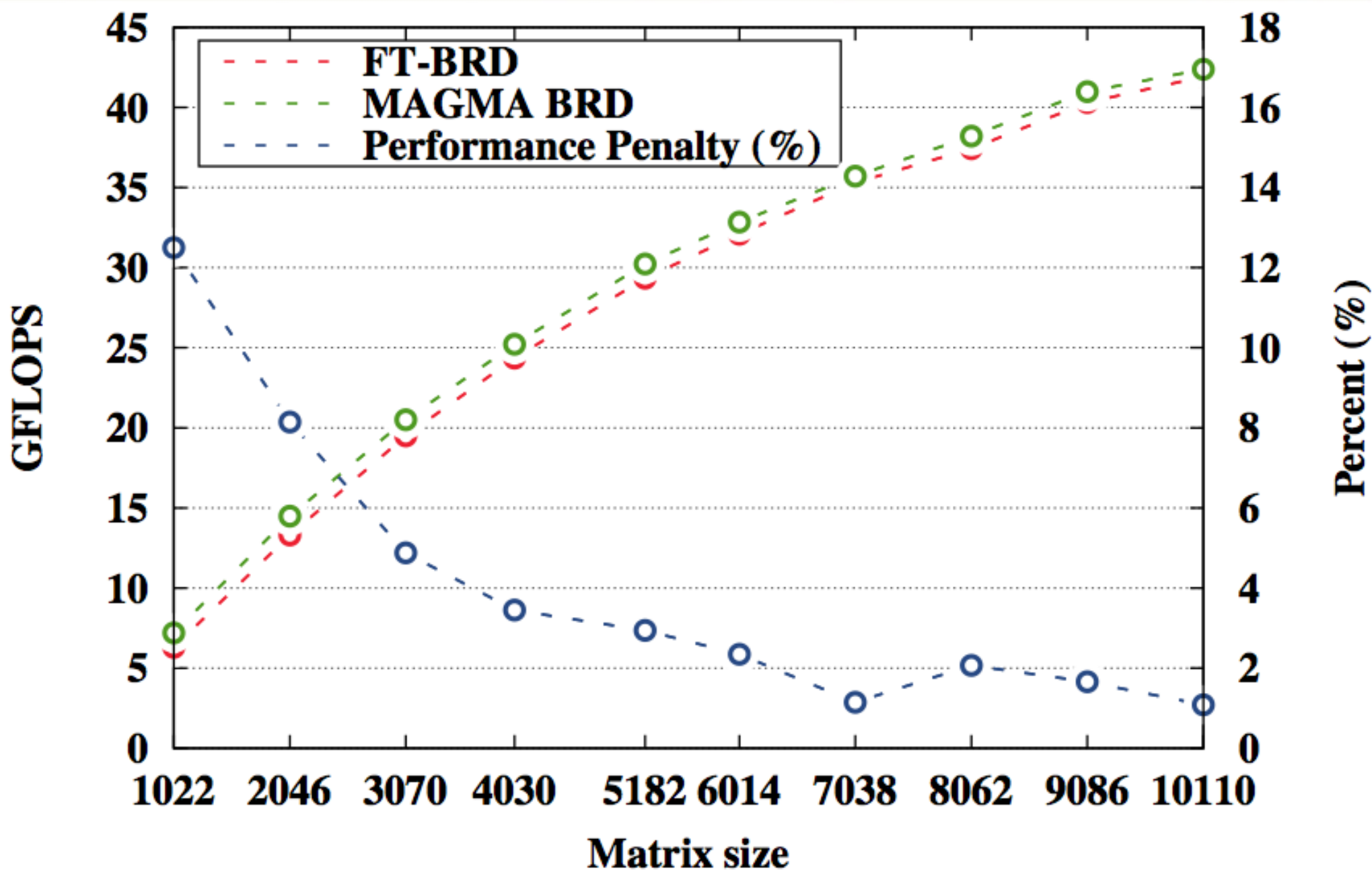
Cores	Process grid	Fault-Tolerant	Classic (ScaLAPACK)
36	6 x 6	5.20×10^{-3}	5.01×10^{-3}
144	12 x 12	3.09×10^{-3}	2.34×10^{-3}
576	24 x 24	2.16×10^{-3}	1.17×10^{-3}
2304	48 x 48	1.36×10^{-3}	6.35×10^{-4}
9216	96 x 96	1.03×10^{-3}	3.37×10^{-4}

$$r_{\infty} = \frac{\|A - QHQ^T\|_{\infty}}{\|A\|_{\infty} N \epsilon}$$

Experiment: BRD on GPU (no error)



Experiment: BRD on GPU (1 error)



Further Details

icl.utk.edu/ft-la



Questions

- What is the cost of error recovery: error early vs. error late?
- 50% of “our” errors are in the same column/row. What to do?
- Don't call it checksum, call it L1-norm regularization
- What is the cost of 2 errors? K errors? Is it K^3 ?
- How does MTBF change if the machine/software can tolerate 1 error?