

Matrix Algebra on GPU and Multicore Architectures

Innovative Computing Laboratory Electrical Engineering and Computer Science University of Tennessee

Piotr Luszczek (presenter)

web.eecs.utk.edu/~luszczek/conf/

MAGMA: LAPACK for GPUs

- **MAGMA**
	- Matrix Algebra for GPU and Multicore Architecture
	- To provide LAPACK/ScaLAPACK on hybrid architectures
	- <http://icl.cs.utk.edu/magma/>
- MAGMA BLAS
	- A subset of BLAS for GPUs
	- Highly optimized for NVIDIA GPGPUs
	- Fast GEMM for Fermi
- MAGMA developers & collaborators
	- UTK, UC Berkeley, UC Denver, INRIA (France), KAUST (Saudi Arabia)

ICLOUT

– Community effort, similar to LAPACK/ScaLAPACK

● *Matlab, Python* ● *Linux, Windows, Mac OS X* ● *C/C++, Fortran*

ICLOUT

MAGMA Functionality

- 80+ hybrid algorithms have been developed (total of 320+ routines)
	- Every algorithm is in 4 precisions $(s/c/d/z)$
	- There are 3 mixed precision algorithms (zc & ds)
	- These are hybrid algorithms, expressed in terms of BLAS
	- MAGMA BLAS
- A subset of GPU BLAS, optimized for Tesla and Fermi GPUs

SINGLE GPU

Hybrid LAPACK algorithms with static scheduling and LAPACK data layout

MULTI-GPU **STATIC**

Hybrid LAPACK algorithms with 1D block cyclic static scheduling and LAPACK data layout

MULTI-GPU **DYNAMIC**

Tile algorithms with StarPU scheduling and tile matrix layout

MAGMA Methodology Overview

A methodology to use all available resources:

- MAGMA uses hybridization methodology based on
	- Representing linear algebra algorithms as collections of tasks and data dependencies among them
	- Properly scheduling tasks' execution over multicore and GPU hardware components
- Successfully applied to fundamental linear algebra algorithms
	- One- and two-sided factorizations and solvers
	- Iterative linear and eigensolvers
- Productivity
	- Use high-level description; low-level hidden with proper abstractions
	- Leverage prior efforts
	- Exceed the performance of homogeneous solutions

Hybrid CPU+GPU Hybrid CPU+GPU algorithms (small tasks for (small tasks for multicores and multicores and large tasks for GPUs) large tasks for GPUs)

Hybrid Algorithms

- Use case: one-sided factorization
	- LU, QR, Cholesky
- Hybridization procedure
	- Panels are factored on CPU using LAPACK (or equivalent)
		- It is slow on the GPU
		- Off-load from GPU to CPU
	- Trailing matrix updates are done on the GPU
	- Look-ahead helps in hiding communication and panel factorization

...

A Hybrid Algorithm Example

• Left-looking hybrid Cholesky factorization in MAGMA

```
for (i=0; i \le n; i += nb) {
 1
 \overline{2}ib = min(nb, n - i):
 3
           magma zherk (MagmaUpper, MagmaConjTrans,
                           jb, j, m_one, dA(0, j), ldda, one, dA(j, j), ldda, queue);
           magma zgetmatrix async(ib. jb. dA(i,j), Idda, work, 0, jb. queue, &event);
 4
 5
           if (i+ib < n)
 6
              magma zgemm(MagmaConjTrans, MagmaNoTrans, jb, n-j-jb, j, mz one,
                                dA(0, i), Idda, dA(0, i+jb), Idda, z one, dA(i, i+jb), Idda, queue);
 \overline{7}magma event sync(event);
 8
           lapackf77 zpotrf(MagmaUpperStr, &jb, work, &jb, info);
 9
           if ( *info := 0 )
10*info += i:
11magma zsetmatrix async(\phijb, \phi), work, 0, \phi), dA(\phi), Idda, queue, & event);
12if (i+jb < n) {
13
              magma event sync(event);
14
              magma ztrsm (MagmaLeft, MagmaUpper, MagmaConjTrans, MagmaNonUnit,
                              jb, n-j-jb, z one, dA(i, i), Idda, dA(i, i+i b), Idda, queue);
          }
```
- \cdot The difference with LAPACK the 4 additional lines in red
- Line 8 (done on CPU) is overlapped with work on the GPU (from line 6)

LU Factorization (single GPU)

GPU Fermi C2050 (448 CUDA Cores @ 1.15 GHz) **AMD Istanbul CPU** + Intel Q9300 (4 cores @ 2.50 GHz) [8 sockets x 6 cores (48 cores) @2.8GHz] DP peak $515 + 40$ GFlop/s DP peak 538 GFlop/s Power * ~ 220 W Power * ~1,022 W

* Computation consumed power rate (total system rate minus idle rate), measured with KILL A WATT PS, Model P430

From Single to Multi-GPU Support

- Data distribution
	- 1-D block-cyclic distribution
- Algorithm
	- GPU holding current panel is sending it to CPU
	- All updates are done in parallel on the GPUs
	- Look-ahead is done with GPU holding the next panel

LU Factorization: Multiple GPUs

Keeneland system, using one node

3 NVIDIA GPUs (M2070 @ 1.1 GHz, 5.4 GB) 2 x 6 Intel Cores (X5660 @ 2.8 GHz, 23 GB)

ICLOUT

Out of GPU Memory Algorithms

- Perform left-looking factorizations on sub-matrices that fit in the GPU memory (using existing algorithms)
- The rest of the matrix stays on the CPU
- Left-looking versions minimize writing on the CPU

A New Generation of DLA Software

Hybrid Algorithms: One-Sided Transformations

- One-Sided Factorizations
	- LU
	- QR, and
	- Cholesky
- Hybridization
	- Panels (Level 2 BLAS) are factored on CPU using LAPACK
	- Trailing matrix updates (Level 3 BLAS) are done on the GPU using "look-ahead"

Hybrid Algorithms: Two-Sided Transformations

• Two-Sided Factorizations

-
-
-
- Hybridization

– Bidiagonal singular values

- Tridiagonal symmetric/generalized eigenvalues
- Upper Hessenberg non-symmetric eigenvalues
- Trailing matrix updates (Level 3 BLAS) are done on the GPU
	- Similar to the one-sided factorizations
- Panels (Level 2 BLAS) are hybrid
	- Operations with memory footprint restricted to the panel are done on CPU
	- The time consuming matrix-vector products involving the entire trailing matrix are done on the GPU

Additional 4x Speedup from Faster GPU BLAS

DSYTRD (symmetric tri-diag Reduction)

Keenland 3 NVIDIA Fermi M2070 1.1 GHz 5.4 GiB; 2x6 Intel Xeon X5660 2.8 GHz 26 GiB

A. Haidar, S. Tomov, J. Dongarra, T. Schulthess, and R. Solca, *A novel hybrid CPU-GPU generalized eigensolver for electronic structure calculations based on fine grained memory aware tasks***, ICL Technical report, 03/2012.**

Multi-GPU Two-Sided Factorizations

• Need HPC multi-GPU Level 2 BLAS (e.g., 50% of flops in the

Performance of DSYMV on M2090's

T. Dong, J. Dongarra, S. Tomov, I. Yamazaki, T. Schulthess, and R. Solca, *Symmetric dense matrix-vector multiplication on multiple GPUs and its application to symmetric dense and sparse eigenvalue problems*, ICL Technical report, 03/2012.

Hybrid Two-Sided Factorizations

Task Splitting & Task Scheduling

From Fast BLAS to Fast Tridiagonalization

Performance of MAGMA DSYTRD on multi M2090 GPUs

- 50 % of the flops are in SYMV
- § Memory bound, i.e. does not scale well on multicore CPUs
- Use the GPU's high memory bandwidth and optimized SYMV
- § 8 x speedup over 12 Intel cores (X5660 @2.8 GHz)

Keeneland system, using one node 3 NVIDIA GPUs (M2070@ 1.1 GHz, 5.4 GB) 2 x 6 Intel Cores (X5660 @ 2.8 GHz, 23 GB)

T. Dong, J. Dongarra, S. Tomov, I. Yamazaki, T. Schulthess, and R. Solca, *Symmetric dense matrix-vector multiplication on multiple GPUs and its application to symmetric dense and sparse eigenvalue problems*, ICL Technical report, 03/2012.

From Static to Dynamic Scheduling …

- Static may stall in situations where work is available
- Hand tuned optimizations
- Hardware heterogeneity
- Kernel heterogeneity
- Separation of concerns
- Dynamic Runtime System

Matrices Over Runtime Systems at Exascale

- MORSE
- Mission statement:
	- "Design dense and sparse linear algebra methods that achieve the fastest possible time to an accurate solution on large-scale Hybrid systems"
- Runtime challenges due to the ever growing hardware complexity
- Algorithmic challenges to exploit the hardware capabilities to the fullest
- Integrated into MAGMA software stack

MAGMA-MORSE: x86 + Multiple GPUs

- Lessons Learned from PLASMA
- New high performance numerical kernels
- StarPU Runtime System
	- Augonnet et. Al, INRIA, Bordeaux
- Use of both: x86 and GPUs leads to Hybrid Computations
- Similar to LAPACK in functionality

High Productivity: Sequential Code

From Sequential Nested-Loop Code to Parallel Execution

```
for (k = 0; k < min(MT, NT); k++) {
   zgeqrt(A[k;k], ...);
   for (n = k+1; n < NT; n++)
      zunmqr(A[k;k], A[k;n], ...);
   for (m = k+1; m < MT; m++) {
      ztsqrt(A[k;k],,A[m;k], ...);
      for (n = k+1; n < NT; n++)
          ztsmqr(A[m;k], A[k;n], A[m;n], ...);
   }
```
}

High Productivity: Parallel Code

From Sequential Nested-Loop Code to Parallel Execution **for (k = 0; k < min(MT, NT); k++) { starPU_Insert_Task(&cl_zgeqrt, A, k, k, ...); for (n = k+1; n < NT; n++) starPU_Insert_Task(&cl_zunmqr(A, k, n, ...); for (m = k+1; m < MT; m++) { starPU_Insert_Task(&cl_ztsqrt(A m, k, ...); for (n = k+1; n < NT; n++) starPU_Insert_Task(&cl_ztsmqr(A, m, n, k, ...); }**

}

Contact Information and Generous Sponsors

Stan Tomov **tomov@eecs.utk.edu**

MAGMA team **http://icl.cs.utk.edu/magma/**

PLASMA team **http://icl.cs.utk.edu/plasma/**

Microsoft[®] $\n *n* _{VDIA.\n}$

Collaborating partners

- University of Tennessee, Knoxville
- University of California, Berkeley
- University of Colorado, Denver
- INRIA, France (StarPU team)
- KAUST, Saudi Arabia

