

Energy-Efficient High-Performance-Computing

Hardware Aware Numerics for Scientific Simulations

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Motivation



Reduce energy consumption!

- Costs over the lifetime of an HPC facility in the range of acquisition costs
- Produces carbon dioxide, a risk for the health and the environment
- Produces heat which reduces hardware reliability

Current state

- Hardware features mechanisms and modes to save energy
- Software (scientific apps) are in general power oblivious



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Motivation



Target scientific application

Sparse linear systems

$$Ax = b$$

arise in many apps. that involve PDEs modeling physical, chemical or economical processes

 Low-cost iterative Krylov-based solvers for large-scale systems: Conjugate Gradient (CG), Preconditioned CG (PCG), GMRES and P-GMRES



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Experimental setup

Hardware platform

- AMD Opteron 6128 (8 cores)@2.0 GHz with 24 GBytes of RAM
- NVIDIA Tesla C1060 (240 cores). Disconnected during CPU-only experiments!
- PCI-Express (16×)





Software implementation of CG, PCG

- AMD: Intel MKL (11.1) for BLAS-1 and own implementation of spmv
- NVIDIA: CUBLAS (3.0) and implementation of spmv based on Garland and Bell's approach
- gcc -03 (4.4.3) and nvcc (3.2)



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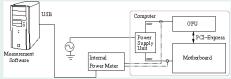
Conclusio

Experimental setup

Measurement setup







ASIC with sampling frequency of 25 Hz



Experimental setup



Linear systems

Matrix name	Size (n)	Nonzeros (nnz)
A318	32,157,432	224,495,280
APACHE2	715,176	4,817,870
AUDIKW_1	943,695	77,651,847
BONES10	914,898	40,878,708
ECOLOGY2	999,999	4,995,991
G3_CIRCUIT	1,585,478	7,660,826
LDOOR	952,203	42,493,817
ND24K	72,000	28,715,634

Solvers Ax = b

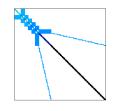
- Iterative: $x_0 \rightarrow x_1 \rightarrow x_2 \rightarrow \cdots \rightarrow x_n \approx x$
- Stopping criterion: $\varepsilon = 10^{-10} ||r_0||_2$
- Initial solution: $x_0 \equiv 0$



Analysis of power consumption

Experiment #1

- Power consumption of CG and PCG on CPU (4 cores) and hybrid CPU (4 cores)+GPU
- G3_CIRCUIT (moderate dimension, complex sparsity pattern)





CG method



Hardware	# iter	Time [s]	Energy consumption [Wh		tion [Wh]
			Chipset	GPU	Total
CPU 4c	21,424	1,076.97	42.18	-	42.18
GPU 4c	21,467	198.43	8.04	3.44	11.48

- Hybrid CPU-GPU code clearly outperforms CPU one in both performance (5×) and energy (4×)
- Energy gap mostly from reduction in execution time:

CPU 4 c	GPU 4 c
$\frac{42.18}{1,076.97} \cdot 3,600 = 140.0 \text{ W}$	$\frac{11.48}{198.43} \cdot 3,600 = 208.2 \text{ W}$



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PCG method (Jacobi preconditioner)

Hardware	# iter	Time [s]	Energy consumption [Wh]		tion [Wh]
			Chipset	GPU	Total
CPU 4c	4,613	348.79	13.31	-	13.31
GPU 4c	4,613	46.28	1.89	0.83	2.72

- Important reduction in #iterations: $21,424 \rightarrow 4,613$
- Time/iteration and energy/iteration not significantly increased (preconditioning this matrix only requires diagonal scaling):

CG GPU 4 c	PCG GPU 4 c
$\frac{198.43}{21,467} = 0.0092 \text{ s/iter}$	$\frac{46.28}{4,613} = 0.0100$ s/iter
$\frac{11.48}{21,467} = 5.34 \cdot 10^{-4} \text{ Wh/iter}$	$\frac{2.72}{4,613} = 5.89 \cdot 10^{-4} \text{ Wh/iter}$



Experiment #2

- In general, for memory-bounded operations a decrease of the processor operation frequency can yield energy savings
 - Memory-bounded or I/O-bounded?
 - Decreasing processor frequency impacts memory latency?
- The sparse matrix-vector product is indeed memory-bounded: 2nnz flops vs. nnz memops
- AMD Opteron 6128: 800 MHz 2.0 GHz
- A318 (large size to match powermeter sampling rate)



CG method



Hardware	Freq.	Time	Power/Energy consumption		
			Chipset	GPU	Total
	[MHz]	[s]	[Avg. W]	[Avg. W]	[Wh]
CPU 4c	2,000	1441.78	123.99	-	49.66
CPU 4c	800	1674.62	108.11	-	50.29
GPU 4c	2,000	253.22	149.04	61.89	14.84
GPU 4c	800	254.25	138.50	61.45	14.12

- For the CPU solver, lowering the processor frequency increases the execution time, which blurs savings in power consumption
- For the hybrid CPU-GPU solver, as the computationally intensive parts are executed on the GPU, lowering the frequency yields some energy savings... Why not larger?



Problems with DVFS



- CPU computations slower when using DVFS
- GPU and CPU operate in asynchronous mode but for GPU kernel calls, the CPU is set into *busy wait* which is very energy-inefficient
- Alternatives:
 - (i) Plain solver(ii) Solver + DVFS during GPU execution



DVFS



DVFS

Power-friendly CPU modes





CG: Energy of chipset+GPU



matrix	Energy	consumption [Wh]	improvement [%]
	(i)	(ii)	(i)→(ii)
A318	14.84	14.12	5.1
APACHE2	1.98	1.99	-0.5
AUDIKW_1	nc	convergence	-
BONES10	nc	o convergence	-
ECOLOGY2	2.30	2.27	-1.3
G3_CIRCUIT	11.48	11.11	3.3
LDOOR	nc	convergence	-
N24K	26.43	25.42	3.97

A moderate gain, in some cases a loss...





PCG: Energy of chipset+GPU



matrix	Energy of	consumption [Wh]	improvement [%]
	(i)	(ii)	(i)→(ii)
A318	14.84	14.12	5.1
APACHE2	1.75	1.76	-0.6
AUDIKW_1	47.98	38.15	5.2
BONES10	157.32	150.16	4.8
ECOLOGY2	2.51	2.45	2.4
G3_CIRCUIT	2.71	2.38	3.0
LDOOR	43.22	41.18	5.0
N24K	34.62	32.97	5.0

Moderate but more consistent gain



Idle-wait



Idle-wait

Experiment #3

- Solution: set the CPU to "sleep" during the execution of the GPU kernels: Execution time of GPU spmv can be measured and accurately adjusted
- Use of nanosleep() function from sys/time.h
- Alternatives:
 - (i) Plain solver
 (ii) Solver + DVFS during GPU execution
 (iii) Solver + DVFS + idle-wait during GPU execution

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Idle-wait



Idle Wait

Power-friendly CPU modes





Idle-wait

CG: Energy of chipset+GPU



matrix	energy	consum	improve	ment [%]	
	(i)	(ii)	(iii)	(i)→(ii)	(i)→(iii)
A318	14.84	14.12	12.18	5.1	21.8
APACHE2	1.98	1.99	1.82	-0.5	8.8
AUDIKW_1	nc	no convergence			-
BONES10	nc	o converg	gence	-	-
ECOLOGY2	2.30	2.27	2.09	-1.3	10.0
G3_CIRCUIT	11.48	11.11	10.10	3.3	13.7
LDOOR	no convergence			-	-
N24K	26.43	25.42	21.17	3.97	24.8



Idle-wait

PCG: Energy of chipset+GPU



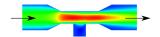
matrix	energy of	consumpt	improve	ment [%]	
	(i)	(ii)	(iii)	(i)→(ii)	(i)→(iii)
A318	14.84	14.12	12.18	5.1	21.8
APACHE2	1.75	1.76	1.64	-0.6	6.7
AUDIKW_1	47.98	45.61	38.15	5.2	25.8
BONES10	157.32	150.16	125.78	4.8	25.1
ECOLOGY2	2.51	2.45	2.29	2.4	9.6
G3_CIRCUIT	2.71	2.63	2.38	3.0	13.9
LDOOR	43.22	41.18	34.79	5.0	24.2
N24K	34.62	32.97	27.64	5.0	25.3



GMRES / Preconditioned GMRES

CFD-application

- Numerical simulation of fluid flow in ventury nozzle
- Navier-Stokes 2D
- Discretized with FEM & linearized
- large, sparse, unsymmetric
 & ill conditioned linear problem



Idle-wait

Example	n	nnz
CFD1	395,009	3,544,321
CFD2	634,453	5,700,633
CFD3	1,019,967	9,182,401





Idle-wait

GMRES / Preconditioned GMRES



Ex.	Solver	Time	Energy
	GMRES	292.23	16.89
CFD1	P-GMRES	194.26	11.30
	gain (%)	33.5	33.1
	GMRES	1104.02	66.66
CFD2	P-GMRES	601.93	36.19
	gain (%)	45.5	45.7
	GMRES	3377.03	231.23
CFD3	P-GMRES	1391.16	86.12
	gain (%)	58.5	62.8

- Strong improvements by adding Jacobi Preconditioner
- Almost linear dependency between time and energy



Idle-wait

Mixed-Iter Variants



Ex.	Solver	Time	Energy		
			Chipset	GPÜ	Total
CFD1	GMRES	292.23	12.00	4.89	16.89
	MPIR GMRES	154.30	6.34	2.71	9.05
	gain (%)	47.2	47.2	44.6	46.4
	P-GMRES	194.26	8.02	3.28	11.30
	MPIR P-GMRES	122.43	5.05	2.30	7.35
	gain (%)	37.0	37.0	29.9	35.0
CFD2	GMRES	1104.02	46.53	20.13	66.66
	MPIR GMRES	640.84	26.89	11.91	38.80
	gain (%)	42.0	42.2	40.8	41.8
	P-GMRES	601.93	25.19	11.00	36.19
	MPIR P-GMRES	416.42	17.47	8.46	25.93
	gain (%)	30.8	30.6	23.1	28.4
CFD3	GMRES	3777.03	160.98	70.25	231.23
	MPIR GMRES	2459.84	104.28	47.74	152.02
	gain (%)	34.9	35.2	32.0	34.2
	P-GMRES	1391.16	59.38	26.74	86.12
	MPIR P-GMRES	1520.79	64.47	28.15	92.62
	gain (%)	- 9.3	-8.6	-5.3	-7.5





Idle-wait

Idle Wait for GMRES



Ex.	Solver	Energy (Wh)		
		Plain	Idle-wait	gain (%)
CFD1	GMRES	16.88	15.65	7.31
	MPIR GMRES	12.38	11.62	8.75
	P-GMRES	9.01	8.22	6.09
	MPIR P-GMRES	7.35	6.61	9.99
CFD2	GMRES	66.66	62.03	6.95
	MPIR GMRES	36.19	33.83	10.22
	P-GMRES	38.79	34.83	6.51
	MPIR P-GMRES	25.94	23.39	9.80
CFD3	GMRES	231.23	217.48	5.95
	MPIR GMRES	86.12	80.88	8.70
	P-GMRES	152.02	138.80	6.08
	MPIR P-GMRES	92.62	84.51	8.75

smaller improvements compared to CG / PCG why?



Idle-wait

Conclusions

Idle Wait for GMRES



Ex.	Solver	Energy (Wh)		
		Plain	Idle-wait	gain (%)
CFD1	GMRES	16.88	15.65	7.31
	MPIR GMRES	12.38	11.62	8.75
	P-GMRES	9.01	8.22	6.09
	MPIR P-GMRES	7.35	6.61	9.99
CFD2	GMRES	66.66	62.03	6.95
	MPIR GMRES	36.19	33.83	10.22
	P-GMRES	38.79	34.83	6.51
	MPIR P-GMRES	25.94	23.39	9.80
CFD3	GMRES	231.23	217.48	5.95
	MPIR GMRES	86.12	80.88	8.70
	P-GMRES	152.02	138.80	6.08
	MPIR P-GMRES	92.62	84.51	8.75

smaller improvements compared to CG / PCG

GPU kernel calls "short" compared to iteration \rightarrow merge computations into one kernel



DV

Conclusions

Conclusions



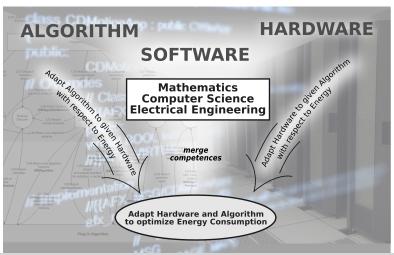
- The concurrency of spmv enables the efficient usage of GPUs, that render important savings in execution time and energy consumption
- For memory-bounded operations, DVFS can potentially render energy savings... but the busy-wait of the host system during the kernel calls still consumes about 80% of full-demand power
- The use of GPU-accelerated HPC-systems combined with power-saving techniques leads to more reduced energy consumption of all test problems without impacting the performance
- Merging more operations into a GPU kernel would trigger longer idle-wait & larger energy savings
- Multi-Iteration kernels would enhance energy efficiency furthermore



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Future Work









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